

양자컴퓨팅 아는 척하기

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목차

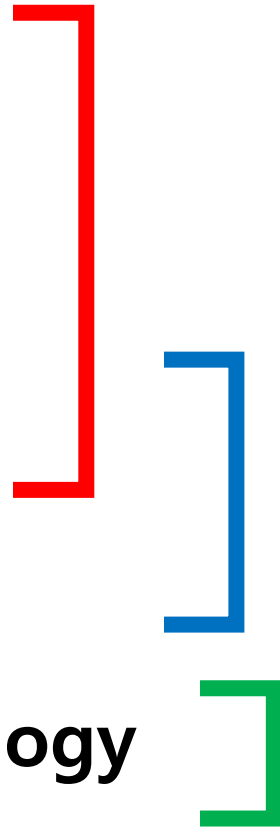
- **Part I. What is "Qubit"?**

- Qubit의 수학적 표현
- 양자 게이트

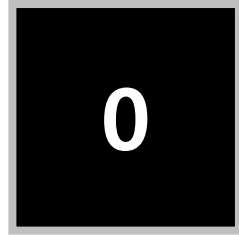
- **Part II. Shor's Algorithm**

- 개요
- 주기탐색 알고리즘

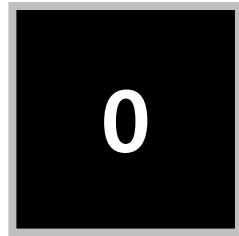
- **Part III. Quantum Cryptology**



Part I. What is "Qubit"?

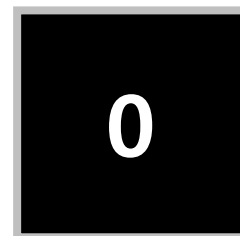


Part I. What is "Qubit"?



=

$a \times$



+ $b \times$



Part I. What is "Qubit"?

$$|\psi\rangle = a|0\rangle + b|1\rangle$$

$$\langle\psi|\psi\rangle = a^2 + b^2 = 1$$

$$|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

$$\text{Bra : } \langle v| = [v_1 \quad v_2 \quad \cdots \quad v_n]$$

$$\text{Ket : } |v\rangle = \begin{bmatrix} v_1 \\ v_2 \\ \vdots \\ v_n \end{bmatrix}$$



Part I. What is “Qubit”?

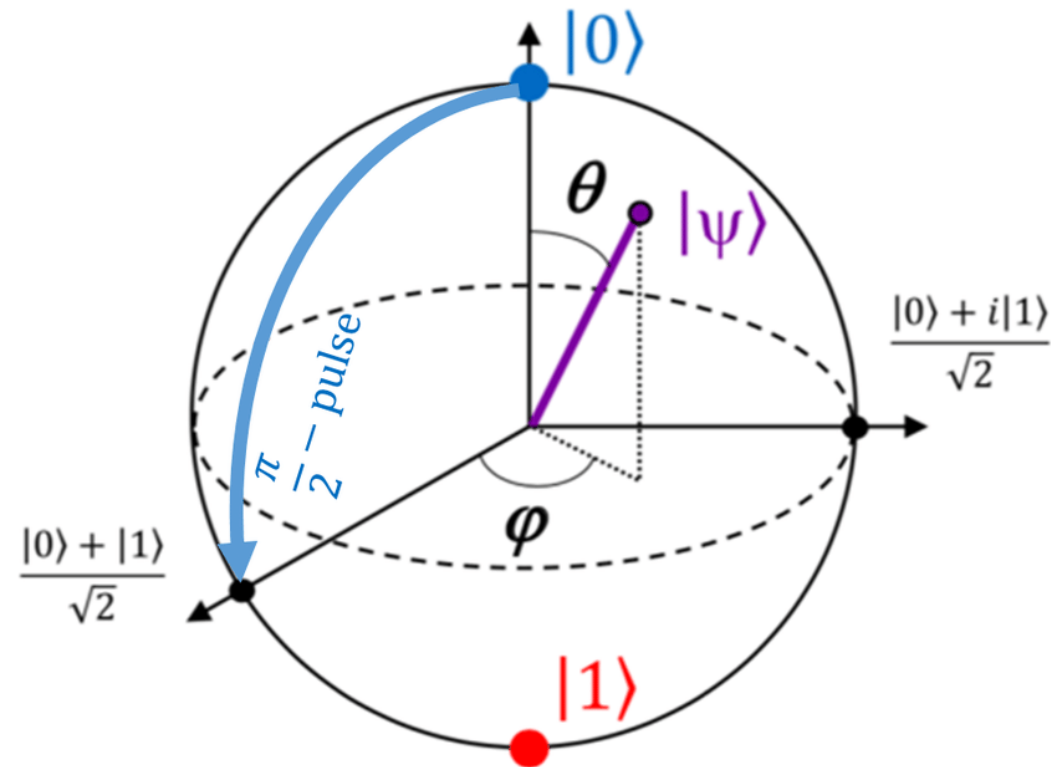
$$|\psi\rangle = a|0\rangle + b|1\rangle = Ae^{i\theta}|0\rangle + Be^{i\phi}|1\rangle$$

$$|\psi\rangle = A|0\rangle + Be^{i\phi}|1\rangle$$

$$A^2 + |Be^{i\phi}|^2 = A^2 + B^2 = 1$$



Part I. What is "Qubit"?



$$|\psi\rangle = \cos \frac{\theta}{2} |0\rangle + e^{i\varphi} \sin \frac{\theta}{2} |1\rangle$$



Part I. What is "Qubit"?

$$\begin{array}{l} |0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\ |1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{array} \quad \longrightarrow \quad |x_1\rangle_1 \otimes |x_2\rangle_2 \cdots$$



Part I. What is "Qubit"?

$$|0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$
$$|1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$



$$|x_1\rangle_1 \otimes |x_2\rangle_2 \cdots$$

Tensor Product



Part I. What is "Qubit"?

$$\begin{aligned} |0\rangle &= \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\ |1\rangle &= \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{aligned}$$



$$|x_1\rangle_1 \otimes |x_2\rangle_2 \cdots$$

Tensor Product

$$A \otimes B = \begin{bmatrix} a_{11}B & a_{12}B & \cdots & a_{1n}B \\ a_{21}B & a_{22}B & \cdots & a_{2n}B \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1}B & a_{n2}B & \cdots & a_{nn}B \end{bmatrix}$$



Part I. What is "Qubit"?

$$\begin{aligned} |0\rangle &= \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\ |1\rangle &= \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{aligned}$$



$$|x_1\rangle_1 \otimes |x_2\rangle_2 \cdots$$



Tensor Product

$$\begin{bmatrix} 1 & 3 \\ 0 & 2 \end{bmatrix} \otimes \begin{bmatrix} 3 \\ 2 \end{bmatrix} = \begin{bmatrix} 3 & 9 \\ 0 & 6 \\ 2 & 6 \\ 0 & 4 \end{bmatrix}$$



Part I. What is "Qubit"?

$$\begin{array}{l} |0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\ |1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{array} \quad \longrightarrow \quad |x_1\rangle_1 \otimes |x_2\rangle_2 \cdots$$

↓
Tensor Product


$|n\rangle$: 위에서부터 n 번째(0-based) element만이 1

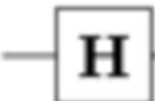


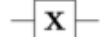

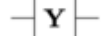
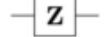
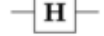
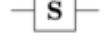
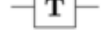

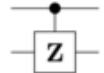



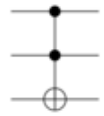
Part I. What is "Qubit"?

Unitary Matrix $UU^* = I$

Pauli-X (X)  $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$

Controlled Not (CNOT, CX)  $\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$

Hadamard (H)  $\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & 1 \end{bmatrix}$

Operator	Gate(s)	Matrix
Pauli-X (X)	 	$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$
Pauli-Y (Y)		$\begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}$
Pauli-Z (Z)		$\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$
Hadamard (H)		$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & 1 \end{bmatrix}$
Phase (S, P)		$\begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$
$\pi/8$ (T)		$\begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix}$
Controlled Not (CNOT, CX)		$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$
Controlled Z (CZ)	 	$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}$
SWAP	 	$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$
Toffoli (CCNOT, CCX, TOFF)		$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$



Part II. Shor's Algorithm

정수 N 의 비자명 약수를 비트 수에 대한 다항 시간 복잡도에 구하는 알고리즘

1. 임의의 정수 $1 < a < N$ 를 선택한다.
 $\gcd(a, N)$ 을 계산하고, 1이 아니라면 $\gcd(a, N)$ 을 반환한다.
2. 함수 $f(x) = a^x \bmod N$ 의 주기 r 을 계산한다.
3. r 이 홀수라면 1로 되돌아가 다시 시행한다.
4. $a^{\frac{r}{2}} \equiv -1 \pmod{N}$ 이라면 1로 되돌아가 다시 시행한다.
5. $\gcd\left(a^{\frac{r}{2}} - 1, N\right), \gcd\left(a^{\frac{r}{2}} + 1, N\right)$ 을 반환한다.



Part II. Shor's Algorithm

$$a^r \equiv 1 \pmod{N}$$

$$\left(a^{\frac{r}{2}}\right)^2 - 1 \equiv 0 \pmod{N}$$

$$\left(a^{\frac{r}{2}} + 1\right)\left(a^{\frac{r}{2}} - 1\right) = mN$$

$$\Rightarrow \gcd\left(a^{\frac{r}{2}} - 1, N\right), \gcd\left(a^{\frac{r}{2}} + 1, N\right)$$



Part II. Shor's Algorithm

$$N^2 < Q = 2^q < 2N^2 \Rightarrow \frac{Q}{r} > N$$

1. q 개의 qubits를 $\frac{1}{\sqrt{Q}} \sum_{x=0}^{Q-1} |x\rangle = \left(\frac{1}{\sqrt{2}} \sum_{x_1=0}^1 |x_1\rangle\right) \otimes \left(\frac{1}{\sqrt{2}} \sum_{x_2=0}^1 |x_2\rangle\right) \otimes \dots$ 로 초기화
Hadamard Gate를 사용하여 수행 가능

2. 함수 f 를 quantum logic gate로 구성(repeated squaring)

$$U_f \frac{1}{\sqrt{Q}} \sum_{x=0}^{Q-1} |x, 0^n\rangle = \frac{1}{\sqrt{Q}} \sum_{x=0}^{Q-1} |x, f(x)\rangle$$

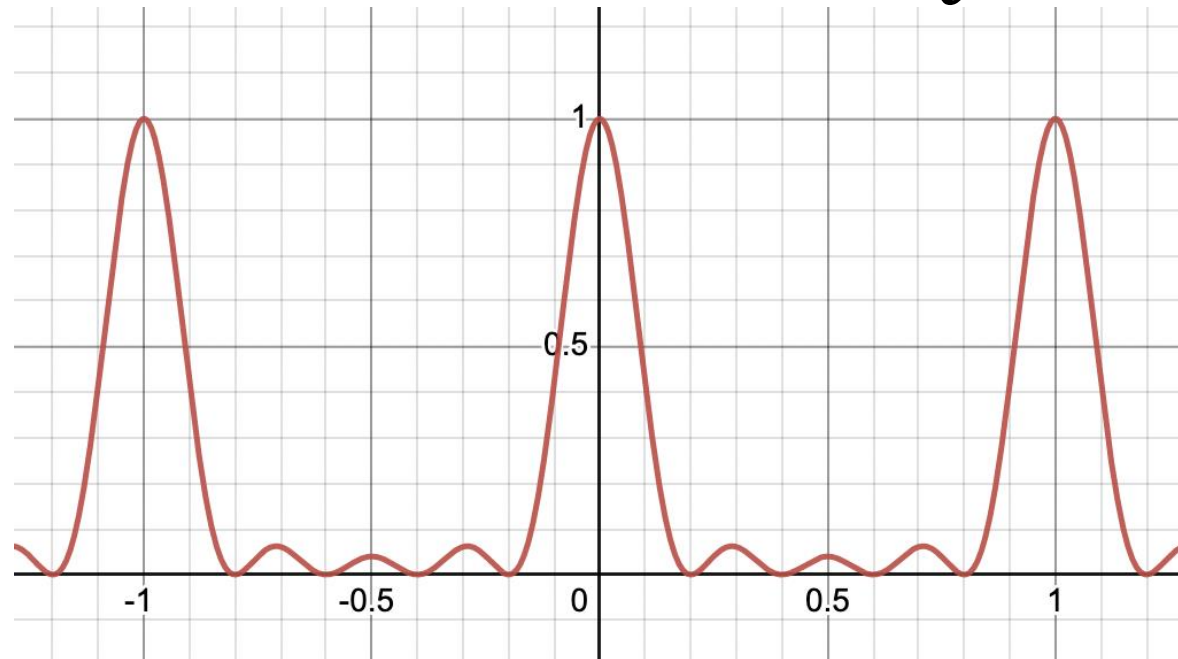
3. 앞의 q 개의 qubits에 QFT를 적용

$$\frac{1}{Q} \sum_{x=0}^{Q-1} \sum_{y=0}^{Q-1} e^{\frac{2\pi ixy}{Q}} |y, f(x)\rangle = \frac{1}{Q} \sum_{z=0}^{Q-1} \sum_{y=0}^{Q-1} \left[\sum_{f(x)=z} e^{\frac{2\pi ixy}{Q}} \right] |y, z\rangle$$



Part II. Shor's Algorithm

$$\Pr(|y, z\rangle) = \frac{1}{Q^2} \frac{\sin^2 \frac{\pi m r y}{Q}}{\sin^2 \frac{\pi r y}{Q}}$$



Part II. Shor's Algorithm

$$\frac{yr}{Q} \approx c \in \mathbb{Z}$$

$$\frac{y}{Q} \approx \frac{c}{r}$$

$$\Rightarrow \frac{d}{s} \approx \frac{c}{r}$$



Part III. Quantum Cryptology

Part II. Shor's Algorithm

정수 N 의 비자명 약수를 비트 수에 대한 다항 시간 복잡도에 구하는 알고리즘

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Part III. Quantum Cryptology

2. 함수 $f(x) = a^x \bmod N$ 의 주기 r 을 계산한다.



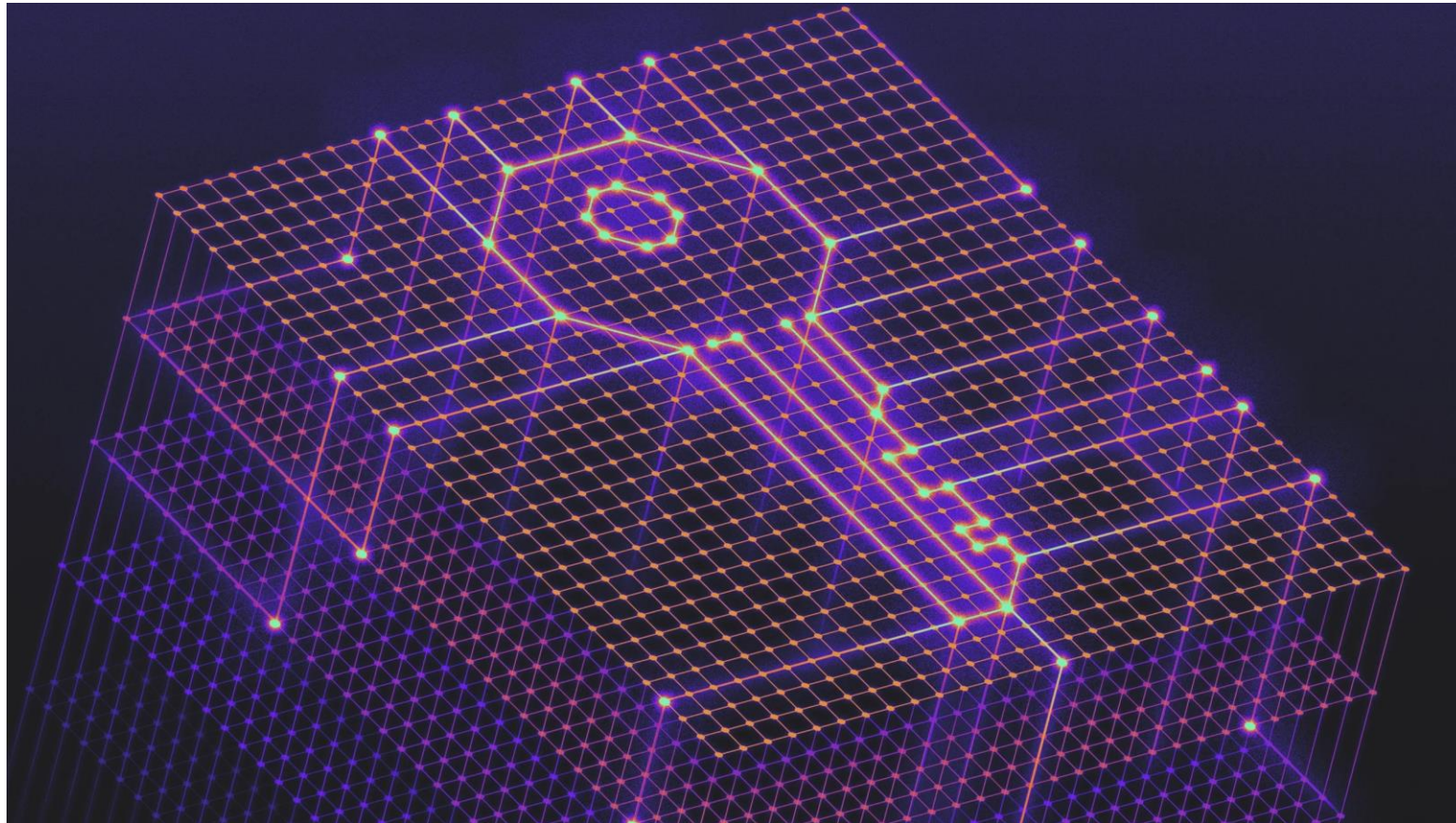
Part III. Quantum Cryptology



Part III. Quantum Cryptology



Part III. Quantum Cryptology





Thank You for Listening

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